

$$\alpha - \alpha R_0^2 = 2R_0^2(1-\alpha)$$

$$\alpha = R_0^2 [2 - 2\alpha + \alpha]$$

$$= R_0^2 [2 - \alpha]$$

$$R_0^2 = \frac{\alpha}{2-\alpha}$$

$$R_0 = \sqrt{\frac{\alpha}{2-\alpha}}$$

Fishy factor of 2.

looks like Petz-Reffy

$Q(z)$ is our $\frac{V(z)}{2}$.

Want R_0 smallest solution

to $R_0 Q'(R_0) = 1$

$$\frac{R_0 V'(R_0)}{2} = 1$$

$$R_0^2 \left[\frac{1-\alpha}{\alpha} + 1 \right] = 1$$
$$= \frac{1}{\alpha}$$

$$\Rightarrow R_0 V'(R_0) = 2$$

$$\Rightarrow R_0^2 = \alpha$$

Now, $V'(r) = \frac{2r(1-\alpha)}{\alpha(1-r^2)}$

$$R_0 = \sqrt{\alpha}$$

\Rightarrow want

$$\frac{2/R_0^2}{(1-R_0^2)} \left(\frac{1-\alpha}{\alpha} \right) = 2$$

$$\Rightarrow R_0^2 \left(\frac{1-\alpha}{\alpha} \right) = 1 - R_0^2$$

good.

$$= \frac{2(1-\alpha)^2}{\alpha^3} \left[\frac{1}{2} \int_0^1 -\frac{w \log w}{(1-w)^3} dw - \frac{1}{2} \log \alpha \left(\frac{\alpha^2}{2(1-\alpha)^2} \right) \right]$$

$$g(x) = \frac{x^2 \log x + x - (x-1)^2 \log(1-x) - 1}{2(x-1)^2}$$

$$g'(x) = \frac{2(x-1)^2 [2x \log x + x - 2(x-1) \log(1-x)] - [x^2 \log x + x - (x-1)^2 \log(1-x) - 1] \cdot 4}{4(x-1)^4}$$

$$= \log x [4x(x-1)^2 - 4x^2(x-1)] + \log(1-x) [-4(x-1)^3 + 4(x-1)^3] + \frac{2(x-1)^2(x+1) - 2(x-1)^2}{-4(x-1)^2}$$

$$= \frac{-8(1-\alpha)^2}{\alpha^3} \left[\frac{w^2 \log w + w - (1-w)^2 \log(1-w) - 1}{2(1-w)^2} \Big|_0^\alpha - \log \alpha \left(\frac{\alpha^2}{2(1-\alpha)^2} \right) \right]$$

$$= \frac{-(1-\alpha)^2}{\alpha^3} \left[\frac{\alpha^2 \log \alpha + \alpha - (1-\alpha)^2 \log(1-\alpha) - 1}{2(1-\alpha)^2} + \frac{1}{2} - \frac{\alpha^2 \log \alpha}{2(1-\alpha)^2} \right]$$

$$= \frac{-(1-\alpha)^2}{\alpha^3} \left[\frac{-(1-\alpha) - (1-\alpha)^2 \log(1-\alpha) + (1-\alpha)^2}{2(1-\alpha)^2} \right]$$

$$\frac{1}{R^2} \log\left(\frac{1}{R}\right) = -\frac{\log R}{R^2} \rightarrow 0$$

$$= \frac{(1-\alpha)}{2\alpha^3} \left[1 + (1-\alpha)^2 \log(1-\alpha) - (1-\alpha) \right]$$

(Scratch Work sample 2/6.)

$$\sum_{j=1}^m \frac{(n-m)_j}{(j-1)!} (1-u)^{j-1} = \sum_{j=0}^{m-1} \frac{(n-m)_{j+1}}{j!} (1-u)^j = (n-m) \sum_{j=0}^{m-1} \frac{(n-m+1)_j}{(1)_j} (1-u)^j$$

$${}_2F_1(a, b, c; z) = \sum_{n=0}^{\infty} \frac{(a)_n (b)_n}{(c)_n} \frac{z^n}{n!}$$

$$= (n-m) \sum_{j=0}^{m-1} \binom{m-1}{j} \frac{(n-m+1)_j (n-j)}{(m-1) \dots (m-j)} (1-u)^j$$

$${}_2F_1(-m, b, c; z) = \sum_{n=0}^m (-1)^n \binom{m}{n} \frac{(b)_n}{(c)_n} z^n$$

$$\sum_{j=0}^{m-1} [n-(m-1)]_j \left[\frac{(1-u)^j}{j!} \right]$$

If $b=c \dots$

$$\sum_{l=m}^{n-(m-1)} n-m+l$$

$$\sum_{l=0}^{m-1} \frac{(n-(m-1))_l}{l!} (\bar{z}\bar{w})^l$$

$$\sum_{l=0}^{n-(m-1)} \binom{n-(m-1)}{l} (\bar{z}\bar{w})^l = (1+\bar{z}\bar{w})^{n-(m-1)}$$

so what? This isn't helpful.

$${}_pF_q(a_1, \dots, a_p; b_1, \dots, b_q; z) = \sum_{n=0}^{\infty} \frac{(a_1)_n \dots (a_p)_n}{(b_1)_n \dots (b_q)_n} \frac{z^n}{n!}$$

$$\sum_{l=0}^{m-1} \frac{(n-(m-1)+l-1)!}{l! (n-(m-1)-1)!} (\bar{z}\bar{w})^l = \sum_{l=0}^{m-1} \frac{n-m+l}{l}$$

(Scratch Work sample 3/6.)

$i \neq q, j \neq l$:

$$\begin{aligned} \frac{1}{\binom{n}{k}} \sum_S \sum_{\substack{j \in S \\ j \neq l}} a_{ij} a_{qj} &= \frac{1}{\binom{n}{k}} \sum_{\substack{r \neq s \\ 1 \leq r, s \leq n}} a_{ir} a_{qs} \binom{n-2}{k-2} \\ &= \frac{k(k-1)}{n(n-1)} \left[\left(\sum_r a_{ir} \right) \left(\sum_s a_{qs} \right) - \sum_r a_{ir} a_{qr} \right] \\ &= \delta_{iq} \end{aligned}$$

$j=l$:

$$\begin{aligned} \frac{1}{\binom{n}{k}} \sum_S \sum_{j \in S} a_{ij} a_{qj} &= \frac{1}{\binom{n}{k}} \sum_{j=1}^n a_{ij} a_{qj} \cdot \binom{n-1}{k-1} \\ &= \frac{k}{n} \delta_{iq} \end{aligned}$$

$$\mathbb{E}[AT_{S'} u_2 - AT_S u \mid S, u]$$

$$\begin{aligned} \mathbb{E}[AT_S(u_2 - u) \mid S, u] &= A \left[\mathbb{E}[T_{S'}(u_2 - u) \mid S, u] \right. \\ &\quad \left. + \mathbb{E}[(T_{S'} - T_S)u \mid S, u] \right] \\ &= -\frac{\Sigma^2}{2k} AT_S u \end{aligned}$$

$$\begin{aligned} &= A \left[-\frac{\Sigma^2}{2k} \mathbb{E}[T_{S'} \mid S] \cdot u \right. \\ &\quad \left. + \left(-\frac{1}{k} T_S u\right) + \left(\frac{1}{n-k}\right)(I - T_S)u \right] \end{aligned}$$

What about pair of S's

$$\begin{aligned} \mathbb{E}[T_{S'} - T_S \mid S] &= \left[\begin{array}{c} \text{pick } z \text{ from } S \\ \downarrow \\ \text{replace w/ } \tau \end{array} \right] = A \left[-\frac{\Sigma^2}{2k} \left(\frac{k-1}{k} T_S + \frac{1}{n-k} (I - T_S) \right) u - \frac{1}{k} T_S u \right. \\ &\quad \left. + \frac{1}{n-k} (I - T_S) u \right] \\ &= A \left[\left(-\frac{\Sigma^2(k-1)}{2k^2} + \frac{\Sigma^2}{2k(n-k)} - \frac{1}{k} + \frac{1}{n-k} \right) T_S u \right. \\ &\quad \left. + \frac{\Sigma^2}{2k(n-k)} u \right] \end{aligned}$$

$$\mathbb{E}[T_{S'} - T_S \mid I_S] = \frac{1}{n-k} \sum_{j \in S} [I - T_S] - (I^u \text{ col})^T \left(\frac{\Sigma^2}{2k(n-k)} + \frac{1}{n-k} \right) u$$

$$\Rightarrow \mathbb{E}[T_{S'} - T_S \mid S] = \frac{1}{n-k} (I - T_S) - \frac{1}{k} T_S \quad \text{not exactly.}$$

(Scratch Work sample 4/6.)

$$\delta_n C_\varepsilon L_\varepsilon \approx \delta_n \left(\frac{2\pi-\theta}{2\pi}\right) n \log(\varepsilon e^{\lambda \delta_n}) = \frac{2\pi-\theta}{2\pi} ?$$

lower bound is

$$\delta_n [A_\varepsilon \cdot L_\varepsilon + B_\varepsilon \cdot L_{1-\varepsilon}]$$

$$\approx \frac{\theta}{2\pi} n \cdot \delta_n \cdot \frac{\lambda}{n \delta_n} = \frac{\theta \lambda}{2\pi}$$

want actually $\delta_n C_\varepsilon L_\varepsilon$ to be small.

$$\approx \delta_n \cdot \left(\frac{n\theta}{2\pi}\right)$$

$$\delta_n \cdot n \cdot L_\varepsilon$$

$$\approx \delta_n \cdot n \cdot \log(1 + \varepsilon(e^{\frac{\lambda}{n\delta_n}} - 1))$$

$$C_\varepsilon = \#\{k: p_k \leq \varepsilon\}$$

$$C_\varepsilon \leq n$$

↓
0 (maybe pretty fast)

So could still have ε s.t. $\varepsilon(e^{\frac{\lambda}{n\delta_n}} - 1)$

blows up

$$\log \varepsilon = -\frac{\lambda}{n\delta_n} + \frac{\varepsilon(n)}{n\delta_n}$$

$$\varepsilon = e^{-\frac{\lambda}{n\delta_n} + \frac{\varepsilon(n)}{n\delta_n}}$$

$$\varepsilon(1-\varepsilon) = \left(e^{\frac{\varepsilon(n)-\lambda}{n\delta_n}} \right) / \left(1 - e^{\frac{\varepsilon(n)-\lambda}{n\delta_n}} \right) \log \varepsilon \approx \left(\log \varepsilon + \frac{\lambda}{n\delta_n} \right) n\delta_n$$

$$\text{need } \log \varepsilon = -\frac{\lambda}{n\delta_n} + x$$

$$\text{s.t. } x n\delta_n \rightarrow 0$$

$$\delta_n A_\varepsilon L_{1-\varepsilon} \leq \frac{\delta_n}{n} \frac{\lambda}{n\delta_n} \cdot \frac{\log n}{\varepsilon(1-\varepsilon)}$$

$$= \lambda \left(\frac{\log n}{n} \right) \left(\frac{1}{\varepsilon(1-\varepsilon)} \right)$$

$$\approx \lambda \frac{\log n}{n} = \lambda e^{\log \log n - \log n + \frac{\varepsilon(n)-\lambda}{n\delta_n}}$$

$$\frac{\varepsilon(n)}{n\delta_n}$$

$$\frac{\varepsilon(n)}{n\delta_n} < \lambda + n\delta_n \log \dots$$

$$\text{need } \frac{\varepsilon(n)-\lambda}{n\delta_n} < \log n \dots (+)$$

$$\frac{\int_t^1 (r^2)^{j-1} (1-r^2)^{n-m-1} r dr}{\int_0^1 (r^2)^{j-1} (1-r^2)^{n-m-1} r dr}$$

$$\int_0^t r^{2j-1} (1-r^2)^{n-m-1} \sum_{k=0}^{n-m-1} \binom{n-m-1}{k} (-1)^k r^{2k} \frac{dr}{t}$$

$$s = \frac{r}{t} \quad ds = \frac{dr}{t}$$

$$= t \int_0^1 (ts)^{2j-1} \sum_{k=0}^{n-m-1} \binom{n-m-1}{k} (-1)^k (ts)^{2k} ds \text{ no use.}$$

$$m^{-\frac{1}{2}(n-m-1)} \log m$$

$$(ts)^{2j-1} (1-ts)^{n-m-1}$$

as long as
 $n-m-1 > 0$, fine!

upper bound?

$$\int_{\mathbb{B}_t^c} (|z|^2)^{j-1} \underbrace{(1-|z|^2)^{n-m-1}}_{\leq t^{n-m-1}} d\mu(z) \leq t^{n-m-1} 2\pi \int_t^1 r^{2j-1} dr$$

$$= 2\pi t^{n-m-1} \left(\frac{1-t^{2j}}{2j} \right)$$

Sum over m : $\sum_{j=1}^m \frac{1}{2j} \sim \frac{1}{2} \log m$ $\sum_{j=1}^m \frac{t^{2j}}{2j} \leq \frac{1}{1-t}$ not so great, actually

wanted $t = \frac{1}{n/m} = \frac{1}{m/n}$

(Scratch Work sample 6/6.)